Verifying Interactive Web Programs *

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Abstract

Web programs are important, increasingly representing the primary public interfaces of commercial organizations. Unfortunately, Web programs also exhibit numerous flaws. In addition to the usual correctness problems faced by software, Web programs must contend with numerous subtle user operations such as clicking the Back button or cloning and submitting a page multiple times. Many existing Web verification tools fail to even consider, much less effectively handle, these operations.

This paper describes a model checker designed to identify errors in Web software. We present a technique for automatically generating novel models of Web programs from their source code; these models include the additional control flow enabled by these user operations. In this technique, we exploit a constraint-based approach to avoid overapproximating this control flow; this approach allows us to evade exploding the size of the model. Further, we present a powerful base property language that permits specification of useful Web properties, along with several property idioms that simplify specification of the most common Web properties. Finally, we discuss the implementation of this model checker and a study of its effectiveness.

1 Introduction

The interactive Web is here to stay. Not only are Web sites generated by programs, but they are increasingly playing the role of "services"—accepting inputs from users, combining these with information in databases, and dynamically computing results. Indeed, from most users' perspectives, a corporation such as Amazon.com or eBay *is* a Web site: the browser is their principal, often only, means of interacting with the organization. As a result, the robustness of Web software has become increasingly important.

Web applications operate in a world of complex user operations. Users can click on the Back button, or clone a window and submit a request from each clone. The Back button forces the computation to resume at a prior interaction point; submitting multiple clones causes computation at the same interaction point to resume multiple times. Worse, these user operations are *silent*: they occur in the browser only, and are not reported to the Web application.

The consequence of these complex and silent user operations is that Web programs manifest numerous subtle errors [11]. For instance, travel sites reserve the wrong flight or hotels. Furthermore, programmers who have not anticipated a sequence of operations are likely neither to develop defensively against it nor to subsequently test for it. User experience demonstrates that even the professionally-developed Web sites of commercially successful companies are not immune to these errors.

User operations are prevalent: (somewhat outdated) studies have shown that the Back button accounts for a significant percentage of user actions [6]. Even attempts to program defensively against them may go awry. A recent New York Times article [22] describes such a situation:

> But when I clicked on the National [car rental] price[...], the site responded with this message: "You may have back-buttoned too far." This was my first experience with "back-button" as a verb. I first translated the phrase as, "You may have pushed the back button too many times." Since that was patently untrue, I decoded its true meaning: "We ran out."

In short, any verification tool for the Web that does not account for user operations is not only incomplete, but potentially even misleading.

In this paper, we present a verification technique for Web software that does account for user operations; this technique includes several significant technical contributions. First, we have designed a Web-aware control-flow analysis that generates a model of a Web program from its source code; this model captures the control flow engendered by user operations. Secondly, we have developed a powerful data-flow-analysis-based property language useful for specifying Web properties, along with several property idioms that simplify specification of the most common Web properties. Finally, we have specialized a model checker with Web domain knowledge for precise verification.

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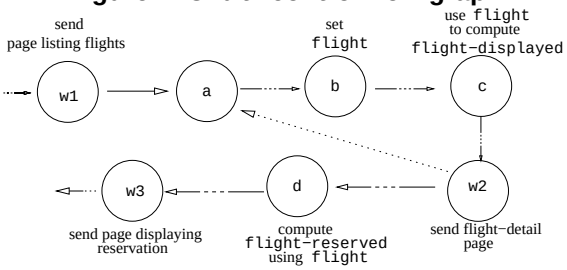
Figure 1. The Orbitz Bug

[Step 1] A user enters the desired dates and destination of his flight; he is then presented with a page listing possible flights, including Flight A and Flight B.
[Step 2] He clicks a link to open the description of Flight A in a new browser window.
[Step 3] Not being particularly enthused about that flight, he returns to the list of flights . . .
[Step 4] and clicks a link to load the description of Flight B, again in a new browser window.
[Step 5] Deciding that Flight A was better after all, he switches back to the window still on the screen showing Flight A . . .
[Step 6] and submits the form, causing a page confirming his reservation to be displayed.
[Result] Orbitz incorrectly makes a reservation on Flight B.

Figure 2. Orbitz control-flow graph



Solid denotes immediate successor; solid-with-ellipsis denotes elided nodes; dotted denotes Web control-flow edge.

2 Motivation and Foundations

We begin by presenting several typical Web program properties that we would like to verify; these properties drive our choice of a particular verification approach.

2.1 Example Properties

A sequence of user operations that exposes an actual bug in the flight-reservation program of Orbitz.com (a travel website) is described in Figure 1. The *Orbitz property* asserts the absence of this bug: the flight described on the page that the user submits in Step 6 (which we will call the `flight-displayed`) should be the same as the actual flight for which his reservation is made (the `flight-reserved`).

One might conclude from this example that all Web sites should have the Orbitz property—that the data used for computation should always correspond to what the user saw on the last page he submitted. However, sometimes it is more desirable to have the *Amazon property*, which is drawn from a desired property of Amazon.com: once the user selects a book to purchase, it should be contained in his shopping cart. In particular, the user should be able to select books in two different browser windows and have both appear in his cart—but this means that the cart will not also satisfy the Orbitz property.

Finally, the *password-page property* prescribes that an authentication page should always be visited before accessing a certain controlled page—starting at page A, you must go through an access-control page B to reach page C.

Each of these three properties involves a notion of temporal sequencing of events. Further, each prescribes that only certain sequences of events should occur on all executions of a system. These qualities suggest the application of model checking.

2.2 Verification Approach: Model Checking

To apply model checking, a developer first creates a model of the system being verified and then writes down the correctness properties with respect to which he would like to verify the model; he then applies a suitable model checking algorithm, which consumes the model and the properties and tells him whether the properties hold for the model.

We now informally work through this methodology for the Orbitz example given above. The desired property is that `flight-reserved` equals `flight-displayed`. Prior work has shown what code for a Web program that exhibits the Orbitz bug would look like [11]; we must extract a model from this code. We could first try the most straightforward technique: from the source code, generate a control-flow graph. A sketch of a control-flow graph is depicted in Figure 2 (ignore the dashed line for now—all the other edges correspond to control-flow from the actual code); though we have elided many details for the sake of presentation, all relevant events (in particular, all assignments to the variable `flight`) are shown. Looking at this model, we notice something interesting: there is no error! Once `flight-displayed` has been computed using the value of the variable `flight`, the program must use the same value compute `flight-reserved`.

Our analysis failed because the sequence of operations listed in Figure 1 that exposed the bug included several instances of the user using his browser to return to and resubmit a previously visited Web page. These actions exploited additional control flow not present in the standard control-flow graph; we must therefore augment our model. The dotted line in Figure 2 shows the control-flow-graph edge necessary to capture the user's ability to return to and resubmit the list-of-flights page (Steps 3 and 4 in Figure 1). (No edge needs to be added corresponding to the browser

window switch in Step 5 because the two pages were generated by the same program expression.) This added control flow exposes the bug—now, we can see that the user might set `flight` to a new value and then submit the page generated using the old one. To develop a sound model of Web programs, we will need to add similar edges for all possible user operations.

2.3 User Operation Calculus

Rather than accounting for each individual operation that a Web browser provides, we use a calculus of primitive user operations due to Graunke, et al. [11]; all traditional browser operations can be expressed in this calculus. Consequently, these primitives are the only user operations about which our verification tool needs to reason.

Like Graunke et al. [11], we distill the Web to a single server and a single user client. The user's client displays Web pages and accepts input; each page includes some text and provides a single form that can be submitted. The client stores a currently active page and a list of previously visited Web pages (which initially contains some start page), and at each step the client can either *submit* the current page's form or *switch* to a previously visited page. When the user *submit*s a form, the server dispatches the request to the correct Web program, which generates and returns a new page based on the client's input. When he *switch*es the client's currently active page, the client does not communicate this change to the server. It is the Web program's lack of knowledge about these *switch*es that causes so many subtle bugs.

In the present work, we make two simplifying assumptions about the user operations: we do not account for the user typing in a URL, and we assume that a Web browser's cache contains all previously visited Web pages. Given these assumptions, we can express the following user operations as combinations of *switch*es and *submit*s: submitting forms, following links, clicking the Back and Forward buttons, choosing a page from the history or bookmarks, refreshing a page, and cloning (i.e., opening a new copy of) a browser window. Verifying using this calculus provides some robustness in the face of new operations that browsers might someday provide—as long as they can be expressed in terms of *switch* and *submit*, we will not need to change our technique. The exact translation of each operation to *switch*es and *submit*s is in our technical report [17].

3 Generating Models From Source

In this section, we describe how we generate a model of the control flow of a Web program from its source code. We assume that the programs we verify are written using special *Web-interaction procedures*. In the present work, we analyze programs written in PLT Scheme that exploit the Web programming procedure `send/suspend` [13]. This procedure consumes a representation of a Web page and sends that page to the user; when he submits the form on the provided page, the Web program resumes computation with the values submitted by the user. In PLT Scheme, this primitive is implemented using continuations (following the lead of Queinnec [20], and as other Web servers also do [1, 2]). When a Web program is written in this form (rather than as a collection of independent scripts), it is not necessary to reason about the marshalling and unmarshalling of data at every Web-interaction point.

We model a Web program P by its *Web control-flow graph* (*WebCFG*). The WebCFG is an augmented control-flow graph (CFG). We define the CFG (N, n_0, E) of a program P as follows:

- The node set N contains one node corresponding to each expression in the source code of P. We denote by $expr(n)$ the source expression corresponding to the CFG node n.
- $n_0 \in N$ is a unique start node
- The edge set $E \subseteq N \times N$ contains an edge (n_1, n_2) iff $expr(n_2)$ might be the next expression to be evaluated after evaluating $expr(n_1)$.

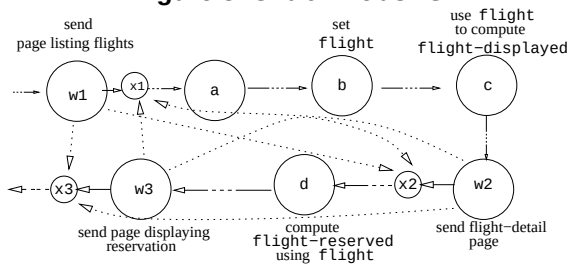
This CFG corresponds to Shivers's 0CFA [21], in that it identifies all executions of a single source expression.

As we saw in the Orbitz example, we must add the control flow engendered by user operations to our model if we are to build a sound verification tool. By the reduction to primitive user operations in Section 2.3, we must only account for *switch* and *submit*. When the user performs a *switch* and then a *submit*, *he causes the program to return from a different Web interaction expression than normal control flow would predict*. That is, when the user performs a *switch*, he changes which Web-interaction call will be returned from on the next *submit*; the program will return from the one that generated the page to which he switched. A *submit* without any *switch*es introduces no new control flow—control proceeds as expected to the successors of the Web interaction expression that generated the submitted page.

To model the control flow enabled by *switch*, we should add a transition from each Web-interaction node w to the successor nodes of each Web-interaction node that the user passed through before reaching w; this corresponds to the user being able to return from (i.e., *switch* to) any previously visited page. Because adding this exact set of edges for all possibles executions will explode the size of the model, we overapproximate by adding to the CFG an edge from each Web interaction node to the successors of every other Web-interaction node (regardless of whether those nodes were passed through on any particular execution).

Formally, we can define the WebCFG $(N^{WebCFG}, n_0^{WebCFG}, E^{WebCFG})$ of a Web program P

Figure 3. Orbitz WebCFG



This figure uses the same arrow convention as Figure 2.

with CFG $(N^{CFG}, n_0^{CFG}, E^{CFG})$ as follows. Let W be the set of Web interaction nodes in P. Let X be a set of fresh nodes, called *post-Web-interaction nodes*, where $|X| = |W|$ and $x_i \in X$ is called the *post-Web-interaction node* corresponding to $w_i \in W$. Then:

- The set N^{WebCFG} of nodes is given by $N^{CFG} \bigcup X$.
- $n_0^{WebCFG} = n_0^{CFG}$.
- The set E^{WebCFG} of edges is the union of
  - $\{(n_1, n_2)$ where $n_1, n_2 \in N^{CFG} - W$ and $(n_1, n_2) \in E^{CFG}\}$
  - $\{(x_i, n)$ where $(w_i, n) \in E^{CFG}\}$
  - $\{(w_i, x_j)$ for all $w_i \in W$ and $x_j \in X\}$.

We could have directly added edges from each Web-interaction node to the successors of every other Web-interaction node, but we have instead introduced the post-Web-interaction nodes to collect these edges—the reasons for this are detailed in Section 6.1.

We construct the WebCFG completely automatically from the source of a Web program using a standard CFG construction technique (Set-Based Analysis [10, 14] approximates the values of procedure-call positions; the CFG can then be constructed by traversing the program's syntax) followed by a simple graph traversal to add the post-Web-interaction nodes and the Web-interaction edges. Figure 3 presents the WebCFG corresponding to the Orbitz CFG considered in Figure 2.

4 Properties

If we annotate the nodes of the WebCFG with elements of some set of atomic propositions, then the graph will describe all traces (atomic proposition sequences) that might occur during execution. The developer formulates a desired property as a set of traces that should occur. Verification then reduces to containment of the former in the latter [23].

Developers specify a set of traces as an automaton whose input alphabet is the set of atomic propositions. Following Naumovich and Clarke [19], whose algorithms we employ, we require developers to write separate automata for safety and liveness properties. (Alpern and Schneider [3] prove that any property described in terms of traces can be decomposed into safety and liveness properties.) Because safety properties are refutable by finite traces, they are expressed as finite-state, finite-word automata with a designated violation (i.e., non-accept) state. Liveness properties are written as deterministic Büchi automata [7]. The determinism is imposed by the model-checking algorithm we have chosen [19], and in principle slightly limits the class of properties we can verify (though we have not encountered this obstacle in practice).

To state concrete properties about Web programs, we must overcome two more challenges:

- To express the Orbitz property, we must be able to talk about strings (the `flight-displayed` and the `flight-reserved`) that appear on Web pages. How can we identify the program expressions that generate these strings?
- The atomic propositions that label WebCFG nodes must be simple enough to be generated automatically, yet rich enough to enable the expression of interesting properties. What should these be?

The next two sections address each of these problems in turn. We then present some example properties, followed by three common property idioms that ease specification.

4.1 Identifying Web-Page Content

How can we associate Web-page contents with the source expressions that generate them? Parsing the HTML fragments in the program to search for strings is likely to be complex, unwieldy and highly sensitive to data and formatting changes. Forcing the developer to use special language constructs to label source expressions is intrusive and not portable. Finally, using static-distance coordinates is brittle in the face of program evolution.

We create a solution that is lightweight, robust in the face of change, and unintrusive by observing that the association is often already in the Web program's source! Web developers often use Cascading Style Sheets (CSS) to tag important page elements with an ID for which independent stylesheets provide formatting directives. We simply ask developers to associate a CSS ID with any Web page element they want to refer to in a property and then to use the ID in an atomic proposition. This ID allows us to identify the source expression that generates the associated Web page element. We similarly use the names of user-input fields to identify source expressions that extract the values of submitted forms. This solution avoids the complexity and brittleness of the other proposed solutions and has the added advantage of being very easy for developers to comprehend—

we have essentially integrated Web presentation elements into the property language.

We refer to a source expression that generates a CSS-tagged HTML element or extracts form input as a *tagged expression*. For example, the Orbitz code might contain tagged expressions that generate HTML with CSS IDs `flight-displayed` and `flight-reserved`, while a search engine might contain one that accesses the user input `query`.

4.2 Atomic Propositions

We give a brief description of our atomic propositions here; a formal treatment is presented in Appendix A.1-2. Some of our atomic propositions are designed for reasoning about misuse of two sets of data bindings: the data local to each page the user sees (e.g., hidden form fields) and the data shared by all pages (e.g., session state or cookies). Many common errors in Web programs result from this misuse [11, 12]. The two data sets have different properties when the user performs browser actions between the generation and the submission of a given Web page: bindings local to each page are guaranteed not to change between page generation and submission, whereas shared bindings may be modified.

Our set of atomic propositions consists of:

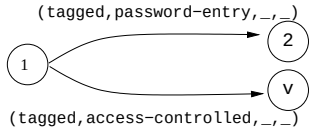
- `tagged` propositions that are true on WebCFG states corresponding to Web program expressions with CSS or user-input tags. These allow the developer to check that certain states have certain values and to reason about the value flow from one expression to another.
- `set` and `join` propositions that are true on states corresponding to operations on shared data. `set` operations replace one value with another, whereas `join` operations add a new value to a collection including all the old ones (for example, mutatively adding to a list is a `join`). These propositions are useful in specifying the Orbitz and Amazon properties.
- `web` and `postweb` propositions that are true at states corresponding to Web-interaction expressions and their successors. These propositions allow the developer to write properties about the sequencing of Web-page generation.

The WebCFG can automatically be annotated with these atomic propositions using the results of a data-flow analysis. However, this analysis requires knowledge about the potential return values of all primitive operations, which means knowing the sets of values that a user might type into each form field. As in the work of Godefroid [5], we presume that the developer has written down some approximation of these values. We assume our tool is given an explicit dictionary-style mapping from field names to values; this mapping could be generated from a more sophisticated (and hence less burdensome to create) user input abstraction such as Godefroid's SmartProfiles.

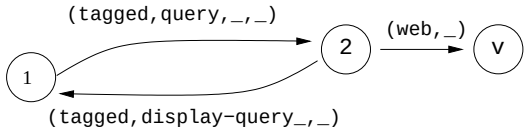
4.3 Example Property Automata

In the following, we label the violation state of a safety property v and the accept states of a liveness property with double concentric circles. Any atomic proposition that is not shown labels a self-loop; any part of a proposition shown with an underscore does not affect atomic proposition matching.

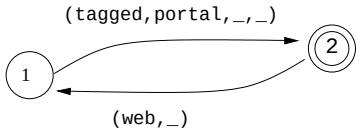
Assuming that there is an expression with CSS tag `password-entry` on the password page and an expression with CSS tag `access-controlled` on the access-controlled page, we can translate the password-page safety property described in Section 2.1 into a property automaton:



Assuming an input tagged `query` and a page element tagged `display-query` that displays the query back to the user on the results page, we can check the property that a search engine always displays the results of a user's query on the next-generated page in part by verifying that the following automaton is satisfied (we also need to verify that `display-query` *takes the value of* `query`; see Section 4.4). This automaton uses the `web` propositions to state that the displayed text is generated before the next page is sent:



Assuming a tag `portal` on a portal page, we can check that a portal page is always eventually reachable by verifying that this automaton is satisfied:



4.4 Property Idioms

Though we could now write the Orbitz and Amazon properties directly as automata, we instead define three property idioms of which they are instances. We write these idioms as automata in Appendix A.3.

So far, we have described the Orbitz property as a relationship between the value of the expression tagged `flight-displayed` and the value of the expression tagged `flight-reserved`: the value of `flight-reserved` must be generated from the value of `flight-displayed` displayed on the page that the user submitted to make his reservation (call this page p_{prev}). This property is implied by the conjunction of two other properties. First, all potential values of `flight-reserved` are also values of `flight-displayed`. Second, no values used in the computation of `flight-reserved` that were present when p_{prev} was generated have changed since its generation. Similarly, we want to capture the Amazon property that once a user selects an item to buy, it it appears in his shopping cart. Assuming appropriate CSS taggings, we state this by saying that the `shopping-cart` contains all values input as `selected-item`.

We offer generalized versions of these properties as idioms in our property language. In the following definitions, let e_1 and e_2 denote expressions in the Web program source.

- We say that e_1 *takes the value of* e_2 iff any potential value of e_1 must also be a value of e_2. The first Orbitz subproperty is an instance of this idiom.
- Let p_{prev} denote the last page the user saw before the evaluation of e_1. Then we say that e_1 is *page* iff no values used in the computation of e_1 that were present when p_{prev} was generated have changed since that page's generation. The second Orbitz subproperty is an instance of this idiom.
- We say that e_1 *accumulates* the values of the e_2 iff the value of e_1 contains the value that e_2 produces at each evaluation (where the exact notion of containment depends on the type of value that e_2 produces). The Amazon property is an instance of this idiom.

5 Verification Process

The model and property language described above are derived from those used in the FLAVERS toolkit [8]: the FLAVERS algorithms consume a model represented as a graph whose nodes are annotated with certain atomic propositions and a property written as an automaton over those same propositions. We may thus reuse the FLAVERS model checking algorithms in our work. We give only an intuitive description of the algorithms; they are presented formally by Naumovich and Clarke [19].

In FLAVERS, a slightly different algorithm is used for verifying liveness properties than for verifying safety properties. Both start with a common subroutine: traverse the model and associate with each model state the property states that are reached at that model state (the atomic propositions reached on model states drive the property automaton). Continue until no model state n can be reached with the property in a state not already associated with n.

Then, to check if a safety property is true for the model, ascertain that no model states are associated with the violation state of the property. Recall that a liveness property is expressed as a deterministic Büchi automaton, and that such an automaton accepts a string iff it reaches an accept state infinitely often. To check a liveness property, form a cross-product graph between each state in the model and the property states that were reached at that state, and prune this graph of all nodes where the property is in an accept state. Then, ascertain that this restricted graph does not contain any strongly connected components. This is a standard model-checking technique (used also in LTL model checking [7, 23]); it relies on the observation that an infinite path where the property never reaches an accept state exists iff such a strongly connected component exists. It is this step that requires the restriction of our property language to deterministic Büchi automata.

At this point, we are able to discover the bug in a model of Orbitz. Using the state labellings from Figure 3 (and only mentioning the depicted nodes), we see that the trace $[w_1, x_1, a, b, c, w_2, x_2, d]$ violates the property that the expression tagged `flight-reserved` is *page*.

6 Improving Precision and Efficiency

In this section, we present two improvements upon our verification technique. The first reduces the number of spurious errors; the second improves the time efficiency of the verification task.

6.1 Constraint Automata for Better Precision

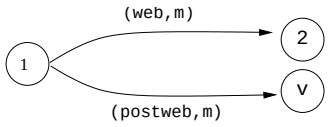
When we discovered the Orbitz bug, we also found that the trace $[w_1, \underline{x_2, d, w_3}, x_1, a, b, c, w_2, x_2, d]$ failed to satisfy the desired property. This corresponds to the user visiting the successors of the second Web-interaction point before he has even gotten to generating the second page, something he clearly cannot actually do. Where did this spurious path come from? When we first defined the Web-CFG, we added edges from each Web-interaction node to the successor of each other Web-interaction node. This overapproximates the control flow introduced by Web interactions: in reality, a user can only *switch* to pages he has seen before, not to any page at all. In this case, the overapproximation resulted in a spurious trace being reported for a program that actually was incorrect. In other cases, it will cause correct programs to be deemed incorrect—for example, the password-page property would never hold, as these

spurious model paths would make it seem as if the user could always jump directly to the access-controlled page.

We can improve the number of correct programs that we deem correct by eliminating these infeasible paths. The naïve way of accomplishing this would be to redesign the WebCFG, adding many slightly augmented copies of the original graph to represent the enabling of new transitions and adding the appropriate transitions between these copies. This approach would cause an exponential explosion in the size of our model, which in turn would drastically increase the time required to check properties over it.

Fortunately, the FLAVERS algorithm gives us a better option. The full FLAVERS algorithm allows the developer to specify any number of *constraint automata* in addition to the property. A constraint automaton, like a safety property, is a finite-state, finite-word automaton with a single violation state that is driven by the propositions reached on the model states. However, the interpretation of reaching the violation state is different: when a constraint automaton is violated, the model path leading to that violation is no longer considered valid; the property is thus allowed to be violated on such paths. The modified FLAVERS kernel can be used directly for the safety property algorithm; it requires only a slight modification (the cross-product graph is now over the model, the property, and all of the constraints) for the liveness-property algorithm. Constraints thus provide an easy and efficient way to prevent certain paths in the model from affecting the results [8].

Helpfully, the constraints needed to remove the spurious paths we introduced in the WebCFG can be generated automatically. We create one constraint of the following form for each Web-interaction node:



Constraint m will be violated on any path where post-Web-interaction node m is visited before Web-interaction node m; since we have a constraint for each Web-interaction node, at least one constraint will be violated on any path that includes a *switch* to a previously unvisited page. The post-Web-interaction nodes are necessary for specifying these constraints: because we have folded all evaluations of an expression into one WebCFG node, any original CFG node can potentially be reached before any given Web interaction node. Thus, none of these original nodes can be used as the atomic proposition that sends the constraint to its violation state without creating the possibility that the constraint will be violated on a valid path.

6.2 Property-language-driven Optimization

Our labelling function associates certain nodes in the WebCFG with the empty set of atomic propositions. Because of the way we have defined the verification process, these unlabeled nodes have no influence on the verification results (since the atomic propositions reached on the model states are the inputs to the property and constraint automata, traversing unlabeled model states will have no effect). Thus, we remove any sequence of unlabeled states from the graph, connecting the predecessors of the sequence directly to its successors. The impact of this optimization on the model size of an example program is presented in Section 8.

7 Soundness and Complexity

We can now state a soundness result of our model checker: if the model checker claims that the WebCFG of a Web program P has a certain property, then that property will hold for all executions of P during which the user performs only *switch*es to previously visited pages and *submit*s. This result follows directly from the soundness of the FLAVERS algorithms and the fact that the WebCFG and the atomic proposition labelling are overapproximations. The WebCFG overapproximates the control flow of the Web program (i.e, if a sequence of expressions is evaluated in order in some execution of the program, then the corresponding sequence of nodes appears in the WebCFG) because the standard CFG construction techniques yield an overapproximation and the edges added to form the WebCFG but not disallowed by the constraints exactly reflect the control flow enabled by the user operations. Our atomic proposition labelling is an overapproximation (i.e., if an atomic proposition holds for a program expression, then the WebCFG node corresponding to that expression is labelled with that proposition) because Set-Based Analysis [10, 14] overapproximates runtime values. These two facts imply that the set of atomic proposition traces along paths in the WebCFG is a superset of the set of atomic proposition traces that actually occur at runtime. By the soundness of FLAVERS, if the model checker claims that a given property holds for a WebCFG, then that property is true for all atomic proposition traces along paths through the WebCFG; in particular, it is true for the subset of those traces that occur at runtime.

The complexity of our method is determined by the complexities of the various phases. The data-flow analysis has worst-case time complexity $O(n^3)$ where there are n expressions in the program source; building the WebCFG then takes time $O(n^2)$, so the time for constructing the model is $O(n^3)$. The FLAVERS safety algorithm takes time $O(n^2 \cdot p \cdot k)$ where p is the number of states in the property and k is the product of the numbers of states in each of the constraints. If a developer uses only the constraints we

generate for the Web control flow, then there will only be one constant-size constraint for each Web-interaction node. Thus, k is $O(w)$, where w denotes the number of Web-interaction nodes, and therefore k is $O(n)$. In this case, we get an overall worst-case upper bound of $O(n^3 \cdot p)$, but we will often do better—w is likely to be much less than n. Checking liveness properties requires additional time for detecting strongly connected components. Our space complexities are the same as those of FLAVERS.

8 Implementation and Results

We have implemented the algorithms described above for constructing the WebCFG and verifying safety and liveness properties. Our implementation accepts Web programs written in PLT Scheme, so we rely on an implementation of Set-Based Analysis called MrFlow [18] for our data-flow analysis. Because our notion of atomic-proposition matching is nuanced (see Appendix A.2), we use a quick reimplementation of the relevant algorithms.

Our implementation makes some simple assumptions to aid in data reasoning. MrFlow does not provide useful value set information about strings, which constitute most Web pages. This is because strings can be combined and decomposed in an arbitrary manner. In contrast, many Web applications do not decompose strings; they only combine strings collected from various sources. (The use of structured forms decreases the need to inspect strings for implicit patterns, such as prefixes that determine gender; this information is instead collected explicitly through separate form fields.) These strings therefore closely resemble collection data structures such as lists (about which MrFlow provides rich value-flow information). We therefore map the string primitives onto list primitives, which enables us to trace the flow of strings through the program. This restriction has sufficed for the programs we have verified.

We have successfully applied our verifier to both correct and incorrect program snippets corresponding to the examples discussed in this paper. Additionally, we have begun to verify CONTINUE [15], a conference-management system that has been used for ISSTA 2004 and many other conferences. Preliminary results are encouraging: an initial WebCFG contained 17,200 nodes, but the property-language-driven state space optimization yielded a model with approximately 500 nodes (the exact number depends on how many expressions are tagged for property use).

9 Related Work

There have been some past efforts to apply formal verification techniques to the Web. De Alfaro [9] uses model checking techniques to verify properties of static (as opposed to interactive and program-generated) Web pages. He treats the page and link structure of the web as a model and then verifies properties written in a slightly restricted μ-calculus over that model. This technique allows him to check many path properties over static Web sites (such as the password-page property) and to present errors as paths through the Web model that violate a given property. Unlike de Alfaro, we are interested in proving properties of interactive, program-generated Web sites.

Godefroid's VeriWeb tool [5] explores interactive Web sites using a special browser that systematically explores all paths up to a specified depth. A user of this tool first makes a model approximating the values that a user might type into the forms of interest. Next, the user specifies properties (such as string containment) about individual Web pages. The verifier then traverses the Web sites of interest and reports errors as sequences of Web operations that lead to a page which violates a property. Like Godefroid, we are concerned with verifying properties of interactive Web sites. However, our work addresses several key limitations of Godefroid's tool. First, Godefroid does not take into account user operations, so his tool is unable to catch errors that only occur in their presence. Our verifier accounts for the control flow enabled by user operations and discovers user-operation-related bugs. Secondly, Godefroid's verification is limited to single-page properties. We provide a method for verifying path properties of interactive Web sites; in particular, we prove properties true of *all paths* through a Web site. We can do this because our verifier operates statically on the program's source, whereas Godefroid's tool is dynamic, effectively "running" the Web site as a Web browser would.

Baresi et al. [4] observe a bug in Amazon and extended UML's OCL with assertions to capture it. These assertions roughly correspond to the properties that we have expressed with *page*; in contrast, we provide a much richer property language. Furthermore, it is unclear how to verify a program against their assertions, as the authors provide neither an algorithm nor a mapping to traditional OCL verifiers.

Our formal model for Web operations is given by Graunke et al. [11]. Using this model, the authors created a type system that statically discovers abuses of the values filled into form fields and devised a strategy for detecting data inconsistency problems such as the Orbitz bug. However, these inconsistency problems are only detected dynamically through changes to the server's run-time system. In contrast, our system is static and can provide guarantees about all possible execution sequences.

10 Future Work

In the future, we would like to perform more case studies to further demonstrate the utility of our tool. We expect that these will help us identify more property idioms, eventually

resulting in a catalog of Web verification patterns.

To verify a larger set of Web applications, we must eventually permit richer reasoning about data. In particular, we must support a broader set of operations on strings (especially string decomposition), as well as arithmetic operations. We expect it will be essential to complement our model checker with a theorem prover: the model checker would output data propositions to the theorem prover, which would determine whether the desired property could be proven from those constraints.

One factor that greatly influences the utility of a model checker is the quality of the error traces it provides when a property is violated. Our technique has the advantage of being able to present traces very intuitively as sequences of Web pages and Web operations that lead to a violation. Indeed, we could potentially even generate this output in the WebVCR format [5] so that a developer could sit back and watch as the error is played out.

We have restricted ourselves to analyzing programs that use the `send/suspend` primitive. Many Web programs, unfortunately, are not written in this style. We conjecture two solutions to this problem. First, based on prior work [12], we conjecture that we can use an "inverse CPS transformation" to convert ordinary CGI programs into a form suitable for our verifier. However, such a tool would have to overcome many engineering obstacles. Secondly, we could treat individual CGI scripts as open features and use techniques [16] for reasoning about their composition.

Currently, we do not address the concurrency issues resulting from multiple simultaneous accesses to a server by different clients (which are different from those resulting from repeated sequential submissions of the same page by a client). Given that many Web sites allow multiple users to interact with the same data, this is an important path for future research. We hope to exploit results on atomicity to reduce the sizes of models involving multiple clients. This process might be abetted by the fact that our current design of the WebCFG includes more knowledge about the particular sequencing of some events than may be necessary.

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A Appendix: Full Property Language

1. Data-flow Analysis

Our property language relies on the results of a standard data-flow analysis. In this work, we use Set-Based Analysis [10, 14], which produces three useful outputs. First, it computes an overapproximation of the runtime values of each expression in the source of the program. Second, it generates a set of flow variables for each expression in the source; this set contains one flow variable for each potential value of that expression. The value set enables reasoning about the values of an expression, whereas the flow variable set enables reasoning about value flow between expressions. Thirdly, it computes a dataflow graph.

2. AP and Labelling

Our atomic propositions rely on a distinction between two sets of bindings, the *environment* (data local to a Web page), and the *store* (data shared by all pages). We assume that a Web program has exactly two syntactically identifiable store operations `set` and `join`. A `set` replaces the value of its first argument with the value of the second, whereas a join adds the value of the second to a collection containing all values previously `join`ed to the first.

$expr(n)$ denotes the source expression corresponding to the WebCFG node n. $Tags$ denotes the set of all tags; $nodes(tag)$ denotes the WebCFG nodes tagged with $tag \in Tags$. V denotes the set of all program values, FV denotes the set of all flow variables, and for $n \in N^{WebCFG}$, $V(n)$ and $FV(n)$ denote the value and flow variable sets corresponding to $expr(n)$. Flow-variable expressions are described by the grammar `FVE ::= s | v | SCv` and `SC` $= \mid \subseteq \mid \not\subseteq \mid \supseteq \mid \not\supseteq$ where $s \subseteq FV$ and v is a flow-variable variable.

AP consists of five kinds of tuples, `tagged`, `set`, `join`, `web`, `postweb`. Their types are as follows: $(\texttt{tagged}, Tags, V, FV)$, $(\texttt{set/join}, Exprs, FVE, FVE)$, and $(\texttt{web/postweb}, \mathbb{Z})$. A labelling function $L : N \to \mathcal{P}(AP)$ associates each node in the WebCFG with a set of atomic propositions that are true at that node. For a given n, $L(n)$ includes:

- $(\texttt{value}, tag, V(n), FV(n))$ iff $n \in nodes(tag)$.
- $(\texttt{set/join}, n_x, FV(n_x), FV(n_v))$ iff $expr(n)$ is an expression that `set`s or `join`s $expr(n_x)$ to $expr(n_v)$.
- $(\texttt{web/postweb}, m)$ iff n is Web-interaction or post-Web-interaction node number m (corresponding Web/post-Web nodes have the same number).

There are a few subtleties in how the developer uses these atomic propositions in properties. When writing `set` and `join` propositions explicitly, the developer must either specify the CFG node by source position, not specify it at all, or use a rule to generate properties for the specific model being checked (we use the second and third approaches when we write *page* and *accumulates* as automata in Figure 4). The use of flow-variable expressions exploits the fact that our verification algorithm is parameterized by the definition of atomic-proposition matching (which determines which transition a property should take when a given model proposition is announced). Other than FVE positions, atomic propositions must exactly match. FVE positions are matched as follows: a literal set in the property matches an identical literal set from the model; a flow-variable variable in the property matches any literal set on the model, and the variable is associated with the literal set in a relation kept by the matching routine; a set-constraint in the property matches a literal set in the model iff the constraint is satisfied for all literal sets related to the constraint's variable. This notion of matching allows us to use value-flow relationships in temporal properties (as we do in *takes the value of* and *accumulates* in Figure 4).

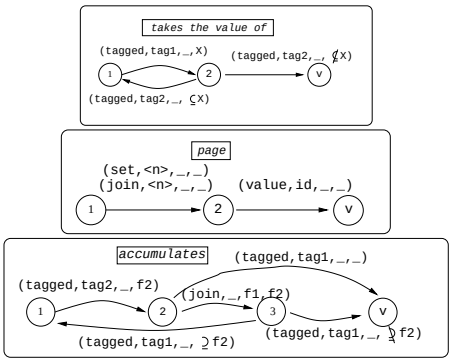

Figure 4. Property Idioms as Automata

$expr(tag2)$ *takes the value of* $expr(tag1)$, $expr(tag)$ is *page*, and $expr(tag1)$ *accumulates* the values of $expr(tag2)$. For *page*, we label the transition from state 1 to state 2 with one `set` and one `join` for each node whose value is used to compute $expr(tag)$ (we identify these using the dataflow graph).

3. Property Idiom Translation

Figure 4 shows how to write our idioms as automata over the full set of atomic propositions. *takes the value of* is simple to express with flow-variable set-constraints. *page* requires that no value used to compute the *page* value can be kept in the store (if some value were kept in the store, the Web control flow would allow it to be mutated between page generation and submission). *accumulates* requires that the accumulated value be `join`ed to the accumulating value.